

Signed Binary Numbers

Unsigned Binary Numbers

We write numbers with as many digits as we need:

0, 99, 65536, 15000, 1979, ...

However, memory locations and CPU registers always hold a constant, fixed number of bits:

- 4 bits: **0000, 1010, 0111, 1101, ...**

- 8 bits: **00000000, 11011001,
00000110, 10111010, ...**

This means there is a biggest *possible* number:

- 4 bits: **$1111_2 = 15$** is the biggest.

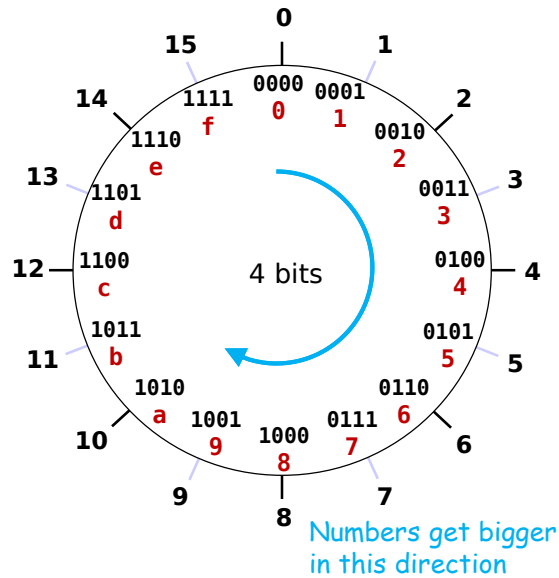
- 8 bits: **$1111\ 1111_2 = 65535$** is the biggest.

Using n bits allows 2^n distinct values in the range:

$\{ 0 \dots 2^n - 1 \}$

4-bit Unsigned Values ...

- $2^4 = 16$ distinct values:
 $\{0, 1, \dots, 15\}$
- Numbers start with $0000b = 0$
- Numbers "wrap around": after $1111b = 15$, back to $0000b = 0$ again



this is
Unsigned Binary

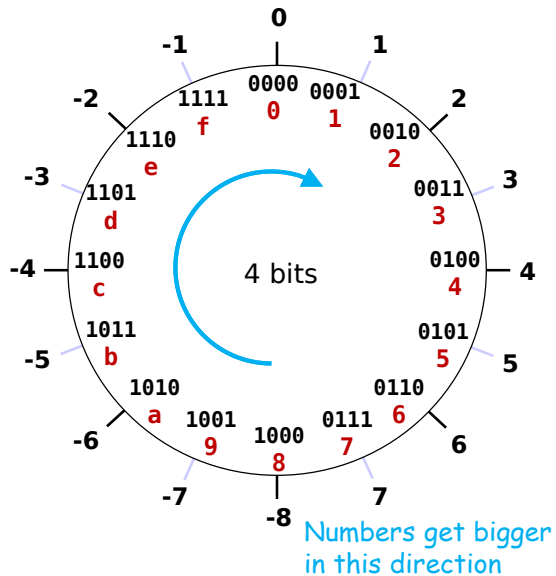
Positive and Negative Numbers

- Subtraction leads to the possibility of numbers smaller than 0:
 $-9 - 7 = 2$, while $7 - 9 = -2$
- How do we represent negative numbers in binary?
 - Fixed number of bits \rightarrow finite set of values
 - 4 bits: **Unsigned** uses this set of 16 values:
 $\{0 \dots 2^4 - 1\} = \{0, 1, \dots, 15\}$
 but this leaves *no room for negatives*
- Use a set of 16 values that includes negatives:
 - 4 bits: use values in the set
 $\{-2^3 \dots 2^3 - 1\} = \{-8, -7, \dots, 7\}$

... 4-bit Signed Values

- Start with **0000** \Leftrightarrow **0**, count in both directions
- Negatives start at **-1**
- Numbers "wrap around" at **-8**

this is
2's Complement



2's-Complement and Hexadecimal

- Hexadecimal – special case
 - Used as "short-hand" representation of binary
- What about 2's-complement (signed) binary?
 - Ignore the sign, just express the binary as hex digits
 - "**Sign-extend**" to get 4 bits per hex digit
- Examples:
 - 01101101 \Leftrightarrow 6d (0x6d in C notation)
 - 11010011 \Leftrightarrow d3 (0xd3 in C notation)
 - 011011 \rightarrow **00**011011 \Leftrightarrow 1b (0x1b in C notation)
 - 11010 \rightarrow **111**11010 \Leftrightarrow fa (0xfa in C notation)

2's Complement Characteristics

- All negative values have a one in the leftmost bit
- All positive values have a zero in the leftmost bit
 - This implies that 0 is a positive value
- Negated zero is still zero (this is a Good Thing!)
- The negative values “go further” than the positive values
 - 4 bits: **-8, -7, -6, -5, -4, -3, -2, -1, 0, +1, +2, +3, +4, +5, +6, +7**
 - 8 bits: **-128, -127, .. 0, .. 126, +127**
 - 32 bits: **-2147483648 .. 0, .. +2147483647**
- *Arithmetic is the same as for unsigned binary!*

Negation

- Negation has to be *invertible*:

$$f(x) = -x$$

$$f(-x) = x$$
- The negation operation must work on the *bit patterns*.
- An algorithm for negations:
 - (1) take the bitwise inverse of the number
 - (2) add 1 to the inverse
 - Symbolically:

$$-x = \bar{x} + 1 \quad (\text{or, } -x = x' + 1)$$

Some examples

$$\begin{array}{r} 7 = 0111 \\ \text{invert: } 1000 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 1001 \end{array}$$

$$\begin{array}{r} -7 = 1001 \\ \text{invert: } 0110 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 0111 \end{array}$$

$$\begin{array}{r} 4 = 0100 \\ \text{invert: } 1011 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 1100 \end{array}$$

$$\begin{array}{r} -4 = 1100 \\ \text{invert: } 0011 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 0100 \end{array}$$

$$\begin{array}{r} 1 = 0001 \\ \text{invert: } 1110 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 1111 \end{array}$$

$$\begin{array}{r} -1 = 1111 \\ \text{invert: } 0000 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 0001 \end{array}$$

$$\begin{array}{r} 0 = 0000 \\ \text{invert: } 1111 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 0000 \end{array}$$

$$\begin{array}{r} -8 = 1000 \\ \text{invert: } 0111 \\ \text{add 1: } + \underline{\quad 1} \\ \hline 1000 \end{array}$$

Oops!

2's-Complement Sign Extension

- Sign extension is the process of extending a value to use more bits
 - The value must remain the same!
- Using more bits in unsigned binary:
 - Just add leading (leftmost) zeroes.
- Using more bits in 2's complement:
 - Copy the leftmost bit into the new positions
 - Positive values stay positive
 - Negative values stay negative

$$\begin{array}{r} +7 = 0000\ 0111 \\ \text{invert: } 1111\ 1000 \\ \text{add 1: } + \underline{\quad 1} \\ \hline -7 = 1111\ 1001 \end{array}$$

$$\begin{array}{r} -7 = \quad \quad 1001 \\ \text{extend to} \\ \text{eight bits:} \\ \hline 1111\ 1001 \end{array}$$

Signed Conversions

Algorithms to convert between
Base-10 and Base-2 signed
numbers

Decimal to Binary

- Convert in two (and a half) steps:
 1. Work with absolute value of base-10 number
 2. Convert absolute value to binary
 - » maintain a leading zero, because this is a positive number!
 3. **IF** the original number was negative,
THEN negate the binary value

$-17 \rightarrow 17 \rightarrow 010001 \rightarrow 101110$
 $\rightarrow 101111$

$17 \rightarrow 17 \rightarrow 010001 \rightarrow (done)$

Decimal to Binary - Code

```

int main(int argc, char **argv)
{
    char bitstring[65], temp[65];
    int value;

    while ( *++argv != NULL) {
        value = atoi( *argv);
        if (value < 0) {
            unsigned2binary_with_leading_0( temp, -(value));
            negate(bitstring, temp);
        } else
            unsigned2binary_with_leading_0( bitstring, value );
        printf("%i --> %s\n", value, bitstring);
    }
}

```

important

Binary to Decimal - 1

- Invert the decimal-to-binary approach:

1. **IF** binary number is negative,
THEN negate to get positive value
2. Convert to decimal as if unsigned
3. **IF** original binary number was negative,
THEN negate the decimal value

011010 → 011010 → 26 → (done)

111010 → 000110 → 6 → -6

This approach is easy for machines/programs

Coding approach for "adding 1"

- Start at the *rightmost* end of the inverted bitstring
 - working left, from least-significant bit to most-significant bit
- Copy each bit to a new string
 - while rightmost bits are "1":
 - a) turn each "1" bit back into a "0" bit
 - b) turn rightmost "0" bit back into a "1" bit
 - copy remainder of bits unchanged
- Remember to add a '\0' bit to the end of the string!

Two Versions of "add_one()" Function

```

/*
 * Two versions of a routine to "add 1" to a bitstring.
 * 1st version uses a while loop, and a switch statement
 * 2nd version uses a for loop, and an if-else if-else statement
 * Both versions preserve any bits other than "0" and "1",
 * so spacing characters are unchanged.
 * Example:
 * "1100 0101" becomes "1100 0110", space character is preserved.
 */
void add_one_whileloop(char *dst, char *dst_end)
{
    while (dst_end >= dst) {
        switch (*dst_end) {
            default:
                break; // do nothing to this char
            case '1':
                *dst_end = '0'; // flip bit back to "0"
                break;
            case '0':
                *dst_end = '1'; // flip bit back to "1"
                return; // ...and quit the routine!
        }
        dst_end--;
    }
}

void add_one_forloop(char *dst, char *dst_end)
{
    for (; dst_end >= dst; dst_end--) {
        if (*dst_end == '1')
            *dst_end = '0';
        else if (*dst_end == '0') {
            *dst_end = '1';
            break; // ...out of for loop
        } else
            ; // leave other char's unchanged
    }
}

```

Other Representations

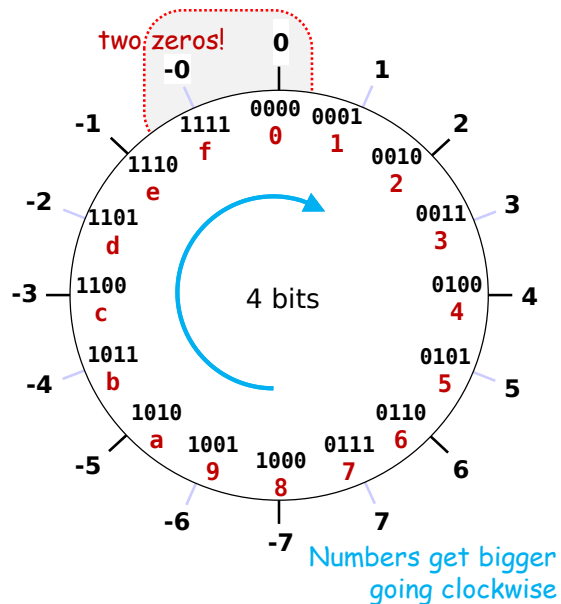
Other ways to represent negative numbers on binary computers

Other ways to make negative numbers:

How about...
just invert
the bits
(without
adding 1)?

Two zeros.
Arithmetic is tricky.
Discontinuity at
 $7 \leftrightarrow -7$

this is
1s' Complement

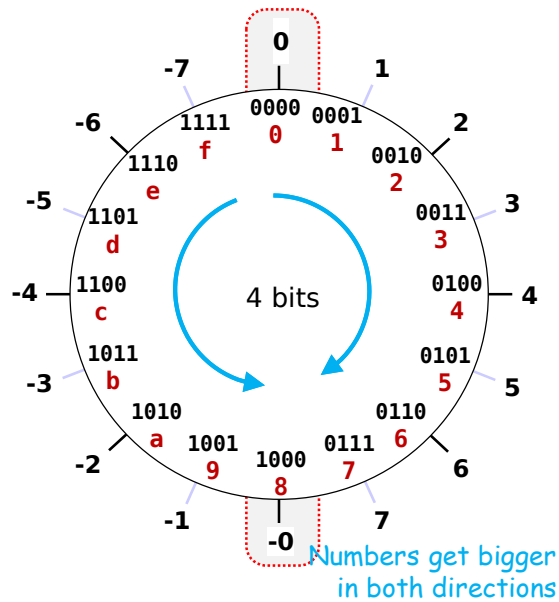


Other ways to make negative numbers:

How about... just set a sign bit?

Two zeros, two discontinuities. Numbers get "bigger" in two directions.

this is **Sign-Magnitude**

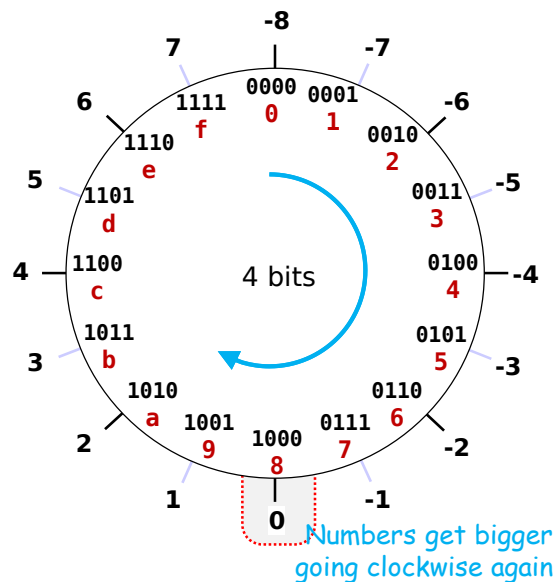


Other ways to make negative numbers:

How about... shift by a constant (such as 8) before converting to binary?

Similar to 2's-complement, but sign bit is reversed.

this is **Excess-8**



Convert Between Decimal and Excess-N

• Decimal → Excess-N:

1. **Add N** to the **decimal** value
2. Convert the result as if for unsigned binary

$$-7 \rightarrow -7 + 8 \rightarrow 1 \rightarrow 0001$$

• Excess-N → Decimal:

1. Convert the binary value as if for unsigned binary
2. **Subtract N** from the result

$$1101 \rightarrow 13 \rightarrow 13 - 8 \rightarrow +5$$

Decimal <--> Excess-8

Summary: Other Representations

• 1s' complement

- Two representations of zero
- Arithmetic works differently than for unsigned binary

• Sign-magnitude

- Two representations of zero
- Not monotonic - arithmetic is very different

• Excess-N

- N can be chosen arbitrarily
- Arithmetic works the same as for unsigned binary
- Single representation of zero
- Negation is more work
- Used for exponents in IEEE-754 floating-point notation

These have been used in the past, but no longer